

Total Neighborhood Number of a Graph

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Abstract

A set S of vertices of a graph G is a total neighborhood set of G if G is the union of the subgraphs induced by the closed neighborhoods of the vertices in S and for every vertex $u \in V$ there exists a vertex $v \in S$ such that u is adjacent to v . The total neighborhood number $n_t(G)$ of G is the minimum cardinality of a total neighborhood set of G . A total neighborhood nomatic partition of G is a partition $\{V_1, V_2, \dots, V_k\}$ of V in which each V_i is a total neighborhood set of G . The total neighborhood nomatic number $n_m(G)$ of G is the maximum order of a partition of the vertex set of G into total neighborhood sets. In this paper, we obtain results about two parameters, the total neighborhood number and total neighborhood nomatic number.

Key words: graph, total neighborhood number, total neighborhood nomatic number.

Mathematics Subject Classification: 05C.

1. Introduction

All graphs considered here are finite, undirected without loops and multiple edges. Let $G=(V, E)$ be a graph with p vertices and q edges. For graph theoretic terminology we refer to Harary³.

A subset D of V is called a dominating set of G if every vertex in $V - D$ is adjacent to

some vertex in D . The domination number $\gamma(G)$ of G is the minimum cardinality of a dominating set of G . A recent survey of $\gamma(G)$ is given in the book by Kulli⁴.

A dominating set D of a graph G without isolated vertices is called a total dominating set of G if the induced subgraph $\langle D \rangle$ has no isolated vertices. The total domination number

$\gamma_t(G)$ of G is the minimum cardinality of a total dominating set of G . This concept was introduced by Cockayne *et al.*¹.

For any vertex $v \in V$, the open neighborhood of V is the set $N(v) = \{u \in V, uv \in E\}$ and the closed neighborhood of v is the set $N[v] = N(v) \cup \{v\}$.

A set S of vertices of G is a neighborhood set if $G = \bigcup_{v \in S} N[v]$. The neighborhood number $n_0(G)$ is the minimum cardinality of a neighborhood set of G , see⁷.

A set S of vertices of G is a total neighborhood set if $G = \bigcup_{v \in S} \langle N[v] \rangle$ and for every vertex u in V , there exists a vertex v in S such that u is adjacent to v . The total neighborhood number $n_t(G)$ of G is the minimum cardinality of a total neighborhood set of G . This concept was introduced by Kulli and Patwari in⁵.

Let $\delta(G)$ denote the minimum degree and $\lceil x \rceil$ ($\lfloor x \rfloor$) the least (greatest) integer greater (less) than or equal to x .

We note that $n_t(G)$ is only defined for G with $\delta(G) \geq 1$.

We need the following result.

*Theorem A*⁶. For a graph \bar{C}_p , $p \geq 5$, $n_0(\bar{C}_p) = 3$.

2. Total Neighborhood Number:

Exact values of $n_t(G)$ for some standard graphs are given below.

Proposition 1. For any complete graph K_p with $p \geq 2$ vertices,

$$n_t(K_p) = 2.$$

Proposition 2. For any complete bipartite graph $K_{m,n}$ with $1 \leq m \leq n$,

$$n_t(K_{m,n}) = m + 1.$$

Proposition 3. For any cycle C_p with $p \geq 3$ vertices,

$$n_t(C_p) = \left\lceil \frac{2p}{3} \right\rceil.$$

Proposition 4. For any path P_p with $p \geq 3$ vertices,

$$n_t(P_p) = \left\lceil \frac{2p}{3} \right\rceil.$$

Proposition 5. For any wheel W_p with $p \geq 4$ vertices,

$$n_t(W_p) = 2.$$

The following results give lower bounds for $n_t(G)$.

Theorem 6. For any graph G without isolated vertices,

$$n_0(G) \leq n_t(G).$$

Theorem 7. For any graph G without isolated vertices,

$$\gamma(G) \leq \gamma_t(G) \leq n_t(G).$$

Theorem 8. If a nontrivial connected

graph G with a γ_t -set S and $V-S$ is independent, then

$$\gamma_t(G) = n_t(G).$$

Proof: Let S be a γ_t -set of G . If $V-S$ is independent, then $G = \bigcup_{u \in S} \langle N[u] \rangle$. Hence S is a n_t -set of G . Thus $n_t(G) \leq \gamma_t(G)$. By Theorem 7, we have $\gamma_t(G) \leq n_t(G)$. Thus $\gamma_t(G) = n_t(G)$.

Remark 9. If G is a tree with $m \geq 2$ cutvertices and every cutvertex is adjacent with at least one endvertex, then

$$n_0(G) = \gamma(G) = \gamma_t(G) = n_t(G).$$

Theorem 10. For any connected graph G with $p \geq 3$ vertices,
 $2 \leq n_t(G) \leq p - 1$.

Furthermore, both inequalities hold if G is C_3 or P_3 .

Theorem 11. Let G be a graph without isolated vertices. Then

$$n_t(G) = p$$

if and only if $G = mK_2, m \geq 1$.

Proof: Suppose $n_t(G) = p$. Let G_1, G_2, \dots, G_m be the components of G . Each $G_i, 1 \leq i \leq m$, has $p_i - 1$ vertices and $G_i \neq K_2$. By Theorem 10,

$$n_t(G_i) \leq p_i - 1, 1 \leq i \leq m.$$

We have

$$n_t(G) = \sum_{i=1}^m n_t(G_i)$$

$$\begin{aligned} &\leq \sum_{i=1}^m (p_i - 1) \\ &\leq p - m \end{aligned}$$

which is a contradiction. Thus $G_i = K_2, 1 \leq i \leq m$ and hence $G = mK_2, m \geq 1$.

Conversely, suppose $G = mK_2, m \geq 1$. Obviously $n_t(G) = p$.

Theorem 12. Let S be a total neighborhood set of G . If $V-S$ is independent, then

$$\alpha_0(G) \leq n_t(G) \leq 2\alpha_0(G).$$

Proof: Suppose S is a total neighborhood set of G and $V-S$ is independent. Then

$$|V-S| \leq \beta_0(G)$$

or $p - n_t(G) \leq \beta_0(G)$.

Thus $p - \beta_0(G) \leq n_t(G)$.

or $\alpha_0(G) \leq n_t(G)$.

We now prove the upper bound. Suppose D is a minimum vertex cover of G . Since G has no isolated vertices, there exists a vertex v_i for every vertex $u_i \in D$ such that $u_i v_i$ is an edge.

Let $S = D \cup \{v_1, v_2, \dots, v_{\alpha_0}\}$.

Then S is a total neighborhood set of G . Thus

$$n_t(G) \leq |S|$$

or $n_t(G) \leq 2\alpha_0(G)$.

Theorem 13. Let G be an r -regular graph with p vertices and $r \geq 1$. Then

$$n_t(G) \leq p - r + 1$$

and this bound is sharp.

Proof: Let $u \in V$. A set $D=(V-N(u)) \cup \{w\}$ is a total neighborhood set of G where $w \in N(u)$. Thus

$$n_t(G) \leq |D| = |(V - N(u)) \cup \{w\}|$$

or $n_t(G) \leq p - r + 1$

The complete graphs $K_p, p \geq 2$, achieve this bound.

Theorem 14. If T is a tree with $m \geq 2$ cutvertices, then

$$n_t(\bar{T}) = 2.$$

Proof: Let u, v be two endvertices of T . Then u, v are adjacent vertices of degree $p - 1$ in \bar{T} . Hence $S = \{u, v\}$ is a total neighborhood set of \bar{T} . Thus

$$n_t(\bar{T}) \leq |S| = 2$$

and since $n_t(G) \geq 2$, we see that

$$n_t(\bar{T}) = 2.$$

Theorem 15. If T is a tree with $m \geq 2$ cutvertices, then

$$n_t(T) + n_t(\bar{T}) \leq p.$$

Furthermore, equality holds if and only if T is P_4 or P_5 or P_6 .

Proof: If T is a tree with $m \geq 2$ cutvertices, then clearly $n_t(T) \leq m$. By Theorem 14, $n_t(\bar{T}) = 2$. Thus

$$n_t(T) + n_t(\bar{T}) \leq m + 2 \leq p.$$

We now prove the second part.

Suppose $n_t(T) + n_t(\bar{T}) = p$. Since $n_t(\bar{T}) = 2$, we have $n_t(T) = p - 2$. Assume T has at least 3 endvertices. Then $n_t(T) \leq p - 3$; a contradiction. Thus T has exactly 2 endvertices. Then T is a path.

Suppose $T = P_p, p \geq 7$. By Proposition 4,

$$n_t(P_p) = \left\lfloor \frac{2p}{3} \right\rfloor \neq p - 2, \text{ which is a contradiction.}$$

Suppose $T = P_2$ or P_3 . Then \bar{T} has isolated vertices and $n_t(\bar{T})$ does not exist.

Hence $G = P_4$ or P_5 or P_6 .

Converse is obvious.

Theorem 16. If C_p is a cycle with $p \geq 7$ vertices, then

$$n_t(\bar{C}_p) = 3.$$

Proof: By Theorem A, $n_0(\bar{C}_p) = 3, p \geq 5$. By Theorem 6, $n_0(\bar{C}_p) \leq n_t(\bar{C}_p)$.

Thus $3 \leq n_t(\bar{C}_p), p \geq 7$. Let u_1, u_2, \dots, u_p be the vertices of $C_p, p \geq 7$. In \bar{C}_p , we see that $\langle N[u_1] \rangle \cup \langle N[u_2] \rangle = \bar{C}_p - \{u_3, u_p\}$. A set $S = \{u_1, u_2, u_r\}$ is a total neighborhood set of \bar{C}_p , where $N[u_r]$ contains u_1, u_2, u_3 and u_p .

Thus $n_t(\bar{C}_p) \leq 3$. Hence $n_t(\bar{C}_p) = 3$.

Theorem 17. If C_p is a cycle with $p \geq 7$ vertices, then

$$n_t(C_p) + n_t(\bar{C}_p) \leq \frac{2p+11}{3}.$$

Proof: By Proposition 3 and Theorem 16, the result follows.

Theorem 18. If both G and \bar{G} are connected graphs with $p \geq 4$ vertices, then

$$n_t(G) + n_t(\bar{G}) \leq 2(p-1)$$

$$n_t(G) \cdot n_t(\bar{G}) \leq (p-1)^2.$$

Furthermore, the bounds are attained if and only if $G = P_4$.

Proof: Suppose G and \bar{G} are connected. By Theorem 10, $n_t(G) \leq p-1$ and $n_t(\bar{G}) \leq p-1$. Hence both inequalities hold.

The second part is easy to prove, so we omit proof.

3. Total Neighborhood Nomatic Number:

The domatic number $d(G)$ of G is the maximum order of a partition of the vertex set of G into dominating sets. This concept was introduced by Cockayne *et. al.*¹. The total domatic number $d_t(G)$ of G is the maximum order of a partition of the vertex set of G into total dominating sets. This concept was

introduced by Cockayne *et al.*¹. In this section, we establish some basic results on the total neighborhood nomatic number of a graph.

In⁵, Kulli and Patwari introduced the concept of total neighborhood nomatic number as follows:

Let G be a graph without isolated vertices. A total neighborhood nomatic partition of G is a partition $\{v_1, v_2, \dots, v_k\}$ of V in which each V_i is a total neighborhood set of G . The total neighborhood nomatic number $n_m(G)$ of G is the maximum order of a partition of the vertex set of G into total neighborhood sets.

The total neighborhood nomatic number of some standard graphs are given below.

Proposition 19. If K_p is a complete graph with $p \geq 2$ vertices, then

$$n_m(K_p) = \left\lfloor \frac{p}{2} \right\rfloor.$$

Proposition 20. If C_p is a cycle with $p \geq 3$ vertices, then

$$n_m(C_p) = 1.$$

Proposition 21. If $K_{r,s}$ is a complete bipartite graph with $1 \leq r \leq s$, then

$$n_m(K_{r,s}) = 1.$$

Proposition 22. If W_p is a wheel with $p \geq 4$ vertices, then

$$n_m(W_p) = 2.$$

It is easy to see the following result.

Proposition 23. If G is a graph without isolated vertices, then

$$n_m(G) \leq \delta(G).$$

Proposition 24. If G is a nontrivial tree, then $n_m(G) = 1$.

Proposition 25. If G is a graph without isolated vertices, then

$$n_t(G) n_m(G) \leq p.$$

Proof: Suppose $\{v_1, v_2, \dots, v_{n_m}\}$ is a total neighborhood nomatic partition of G . Then $|V_i| \geq n_t(G)$ for each i .

$$\text{Hence } \frac{p}{n_m(G)} \geq n_t(G).$$

$$\text{Thus } n_t(G) n_m(G) \leq p.$$

From Proposition 23 and Proposition 25, we have

Proposition 26. If G is a graph without isolated vertices, then

$$n_m(G) \leq \min \left\{ \delta(G), \frac{p}{n_t(G)} \right\}.$$

Proposition 27. If G is a graph without isolated vertices, then

$$n_m(G) \leq d_t(G).$$

Proof: By Theorem 7, $\gamma_t(G) \leq n_t(G)$. Hence $n_m(G) \leq d_t(G)$.

Theorem 28. If G is a graph without isolated vertices, then

$$n_t(G) + n_m(G) \leq p + \delta(G).$$

Furthermore, equality holds if and only if $G = mK_2, m \geq 1$.

Proof: By Theorem 11 and Proposition 22, the inequality follows:

We now prove the second part.

Suppose $n_t(G) + n_m(G) = p + \delta(G)$. Assume $n_t(G) < p$. Then $n_t(G) > \delta(G)$, which is a contradiction. Thus $n_t(G) = p$ and hence by Theorem 11, $G = mK_2, m \geq 1$.

Converse is obvious.

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